Dynamic MapReduce Clusters on Demand

5th Workshop on Many-Task Computing on Grids and Supercomputers Salt Lake City, Utah

Bogdan Ghit, Nezih Yigitbasi, and Dick Epema

Parallel and Distributed Systems Group

Delft University of Technology

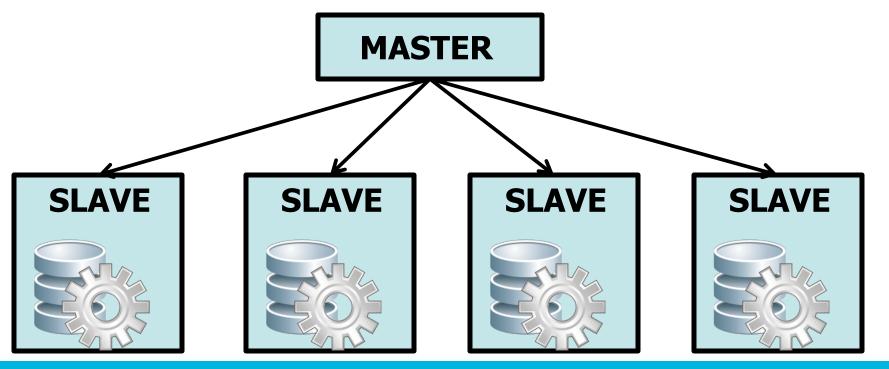
Delft, The Netherlands



MapReduce Overview

- MR cluster
 - ➤ Large-scale data processing
 - Master-slave paradigm

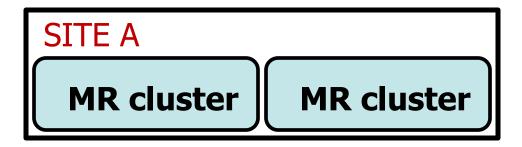
- Components
 - Distributed file system (storage)
 - MapReduce framework (processing)



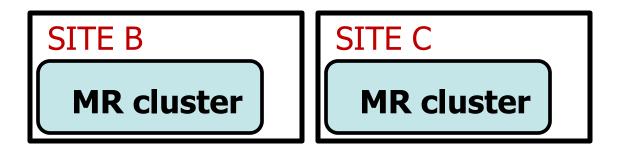


Why Multiple MapReduce Clusters?

Intra-cluster Isolation



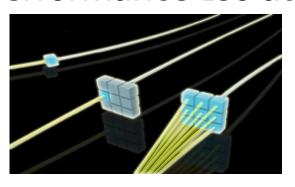
Inter-cluster Isolation





Types of Isolation

Performance Isolation
 Failure Isolation

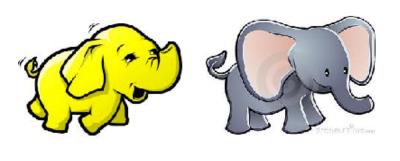




Data Isolation



Version Isolation



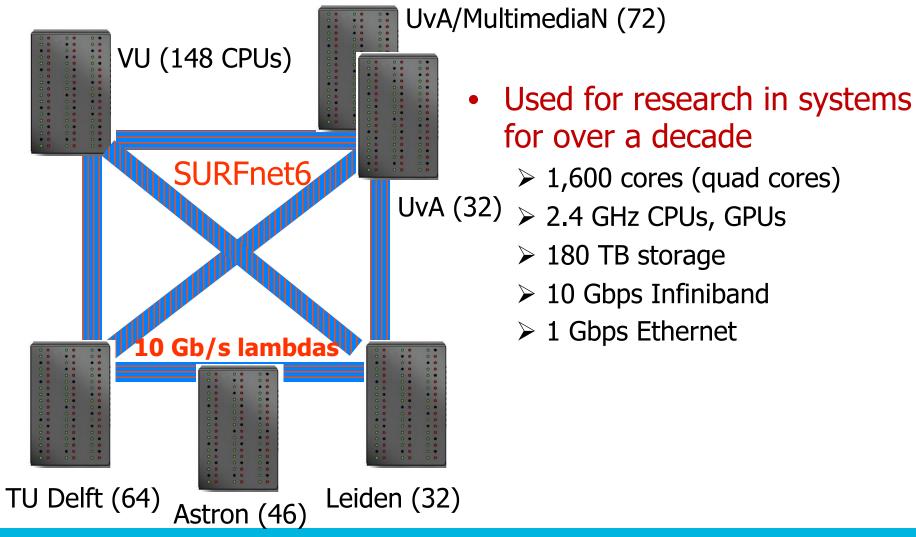
Why Dynamic MapReduce Clusters?

- Improve resource utilization
 - Grow when the workload is too heavy
 - Shrink when resources are idle
- Fairness across multiple MR clusters
 - Redistribute idle resources
 - Allocate resources for new MR clusters





The DAS-4 Infrastructure





Koala Grid Scheduler

- Deployed on DAS-4
- Meta-scheduler, transparent for local schedulers
- · Research vehicle in grid and cloud computing

Features:

- Resource co-allocation
- Scheduling policies
- Various application types



Current runners:

- CSRunner: cycle scavenging apps.
- ➤ OMRunner: co-allocated OpenMPI apps.
- Wrunner: co-allocated workflows
- MR-Runner: MapReduce clusters



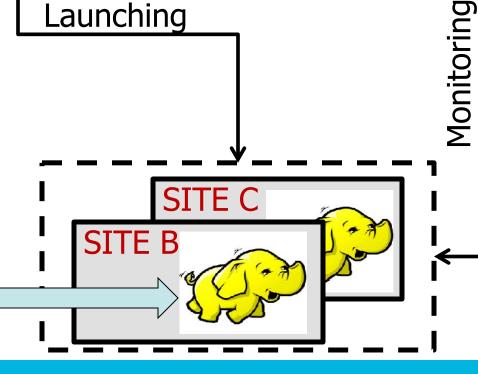
KOALA and MapReduce



Launching

- Users submit jobs to deploy MR clusters
- Koala
 - Schedules MR clusters
 - Stores their meta-data
- MR-Runner
 - > Installs the MR cluster
 - MR job submissions are transparent to Koala

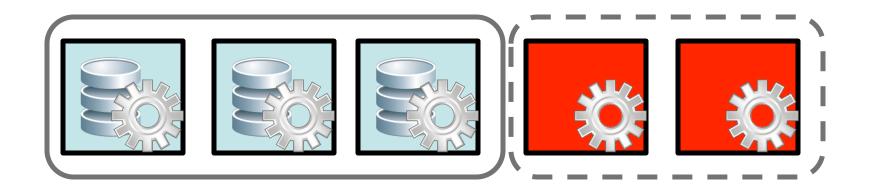
MR jobs





System Model

- Two types of nodes
 - Core nodes: TaskTracker and DataNode
 - Transient nodes: only TaskTracker



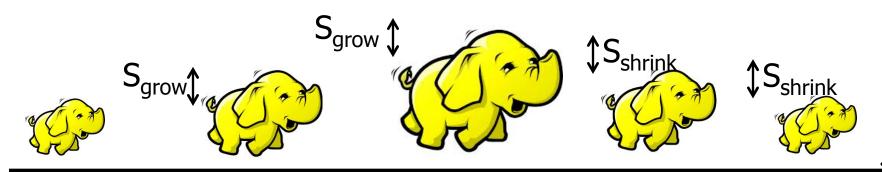


Resizing Mechanism

- Two-level provisioning
 - Koala makes resource offers / reclaims
 - ➤ MR-Runners accept / reject request
- Grow-Shrink Policy (GSP)
 - > MR cluster utilization:

$$F_{\min} \leq \frac{totalTasks}{availSlots} \leq F_{\max}$$

 \succ Size of grow and shrink steps: S_{grow} and S_{shrink}

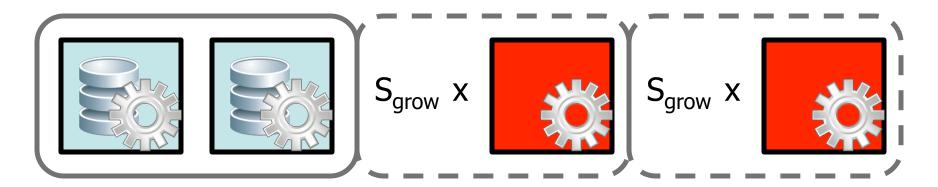


Timeline

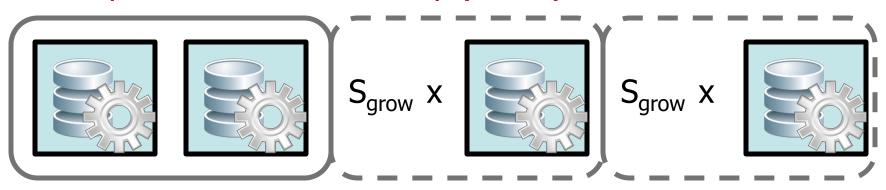


Baseline Policies

• Greedy-Grow Policy (GGP):



Greedy-Grow-with-Data Policy (GGDP):





Setup

- 98% of jobs @ Facebook take less than a minute
- Google reported computations with TB of data
- Two applications: Wordcount and Sort

Workload 1

- Single job
- 100 GB
- Makespan

Workload 2

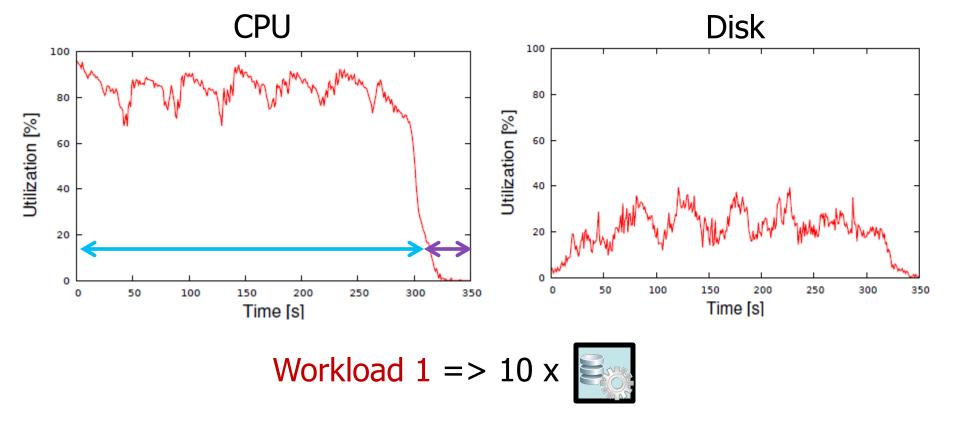
- Single job
- 40 GB, 50 GB
- Makespan

Workload 3

- Stream of 50 jobs
- 1 GB \rightarrow 50 GB
- Average job execution time



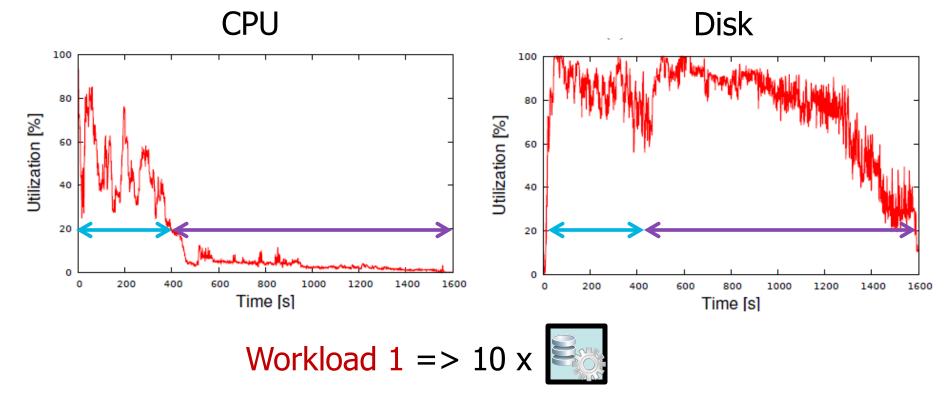
Wordcount



- Wordcount is CPU-bound in the map phase
- Short reduce phase with low CPU utilization



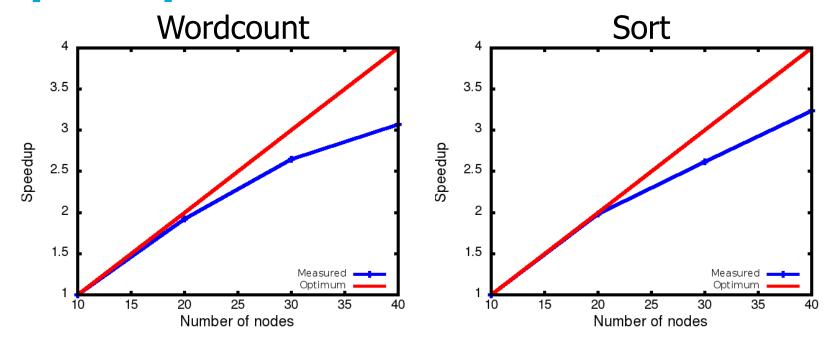
Sort



- Short map phase with 40%-60% CPU utilization
- Long reduce phase which is highly disk intensive



Speedup



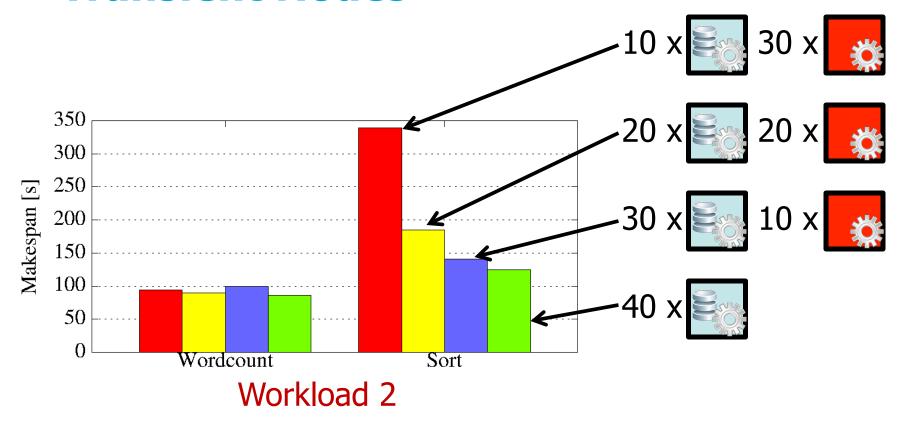
Workload $1 => \{10,20,30,40\} x$



- Speedup relative to an MR cluster with 10 core nodes
- Close to linear speedup on core nodes



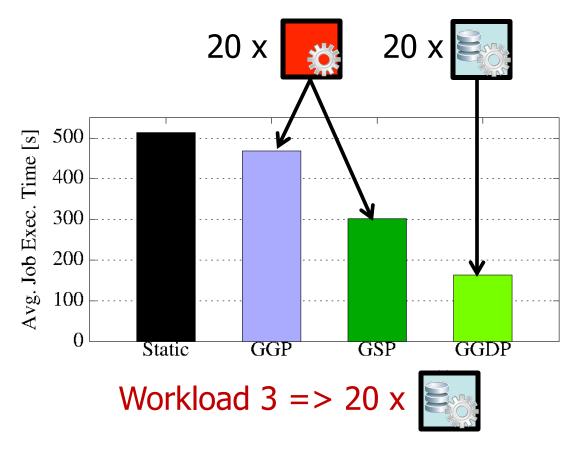
Transient Nodes



Wordcount scales better than Sort on transient nodes



Resizing Performance



Resizing bounds

$$F_{min} = 0.25$$

$$F_{max} = 1.25$$

Resizing steps

$$S_{grow} = 5$$

$$S_{shrink} = 2$$

$$>$$
 GG(D)P

$$S_{grow} = 2$$



Conclusions

- MR clusters on demand
 - System deployed on DAS-4
 - Resizing mechanism

- Performance evaluation
 - > Single jobs workloads
 - Stream of jobs workload

- Distinct applications behave differently with transient nodes
- GSP reduces the job average execution time
- Future Work
 - More policies, more thorough parameter analysis



More Information

- Team: D. Epema, A. Iosup, N. Yigitbasi, S. Shen, Y. Guo, ...
- PDS publication database
 - <u>www.pds.ewi.tudelft.nl/research-publications/publications</u>
- Home pages
 - <u>www.pds.ewi.tudelft.nl/epema</u>
 - www.pds.ewi.tudelft.nl/~iosup
 - www.pds.ewi.tudeltf.nl/ghit

- Web sites:
 - KOALA: www.st.ewi.tudelft.nl/koala

